



The RuleML Family of Web Rule Languages

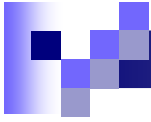
PPSWR'06, Budva, Montenegro, 10 June 2006
Revised, RuleML'06, Athens, GA, 11 Nov. 2006

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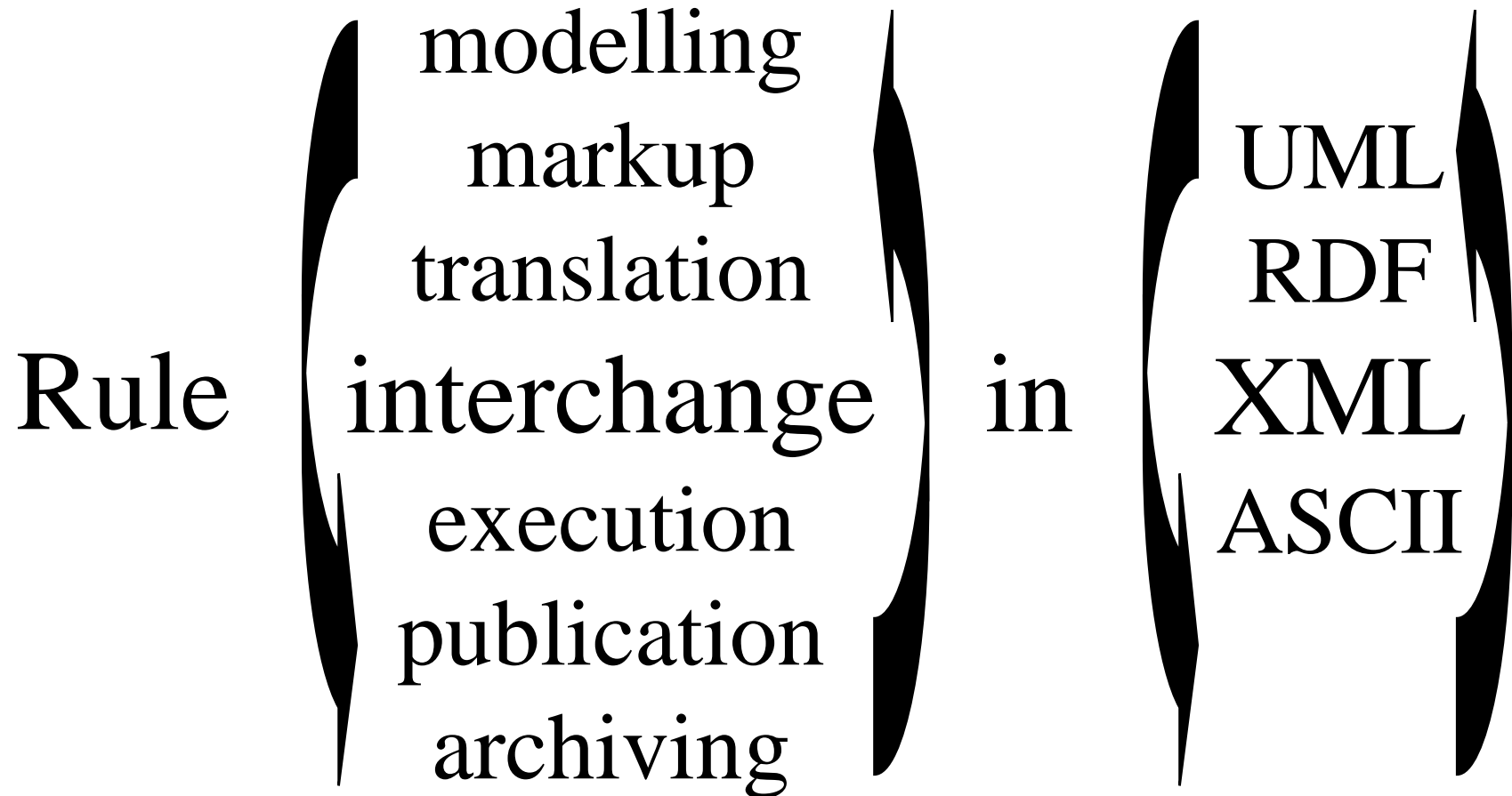


Introduction

- Rules are central to the Semantic Web
- Rule interchange in an open format is important for e-Business
- RuleML is the de facto open language standard for rule interchange/markup
- Collaborating with W3C ([RIF](#)), OMG (PRR, SBVR), OASIS, DARPA-DAML, EU-REWERSE, and other standards/gov'nt bodies



RuleML Enables ...





RuleML Identifies ...

- Expressive sublanguages
 - for Web rules
 - started with
 - *Derivation* rules: extend SQL views
 - *Reaction* rules: extend SQL triggers
 - to empower their subcommunities

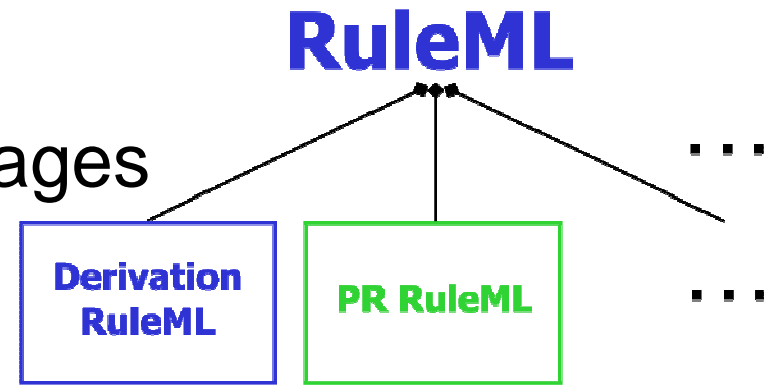


RuleML Specifies ...

- Derivation rules via XML Schema:
 - All sublanguages:(OO) RuleML 0.91
 - First Order Logic: FOL RuleML 0.91
 - With Ontology language: SWRL 0.7
 - A Semantic Web Rule Language
Combining OWL (W3C) and RuleML
 - With Web Services language: SWSL 0.9
- Translators in & out (e.g. Jess) via XSLT

Modular Schemas

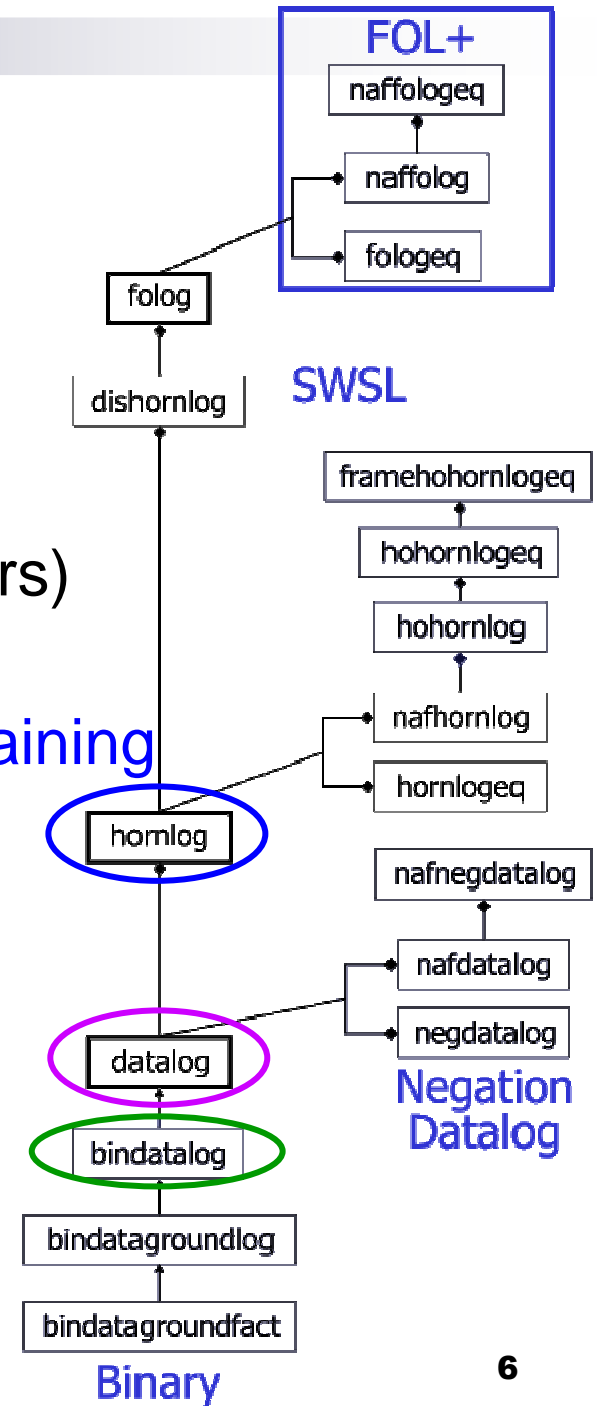
“RuleML is a **family** of sublanguages whose **root** allows access to the language as a whole and whose **members** allow to identify customized subsets of the language.”



- RuleML: Rule Markup Language
 - RuleML derivation rules (shown here) and production rules defined in XML Schema Definition (XSD)
 - Each XSD of the family corresponds to the expressive class of a specific RuleML sublanguage
- The most recent schema specification of RuleML is always available at <http://www.ruleml.org/spec>
- Current release: RuleML 0.91
- Previews: http://wiki.ruleml.org/XSD_Workplan

Schema Modularization

- XSD URIs identify expressive classes
 - Receivers of a rulebase can validate applicability of tools (such as **Datalog** vs. **Hornlog** interpreters)
 - Associated with semantic classes (such as **function-free** vs. **function-containing** Herbrand models)
- Modularization (Official Model)
 - Aggregation: e.g., **Datalog** *part of* **Hornlog**
 - Generalization: e.g., **Bindatalog** *is a* **Datalog**



- **Rectangles** are sublanguages

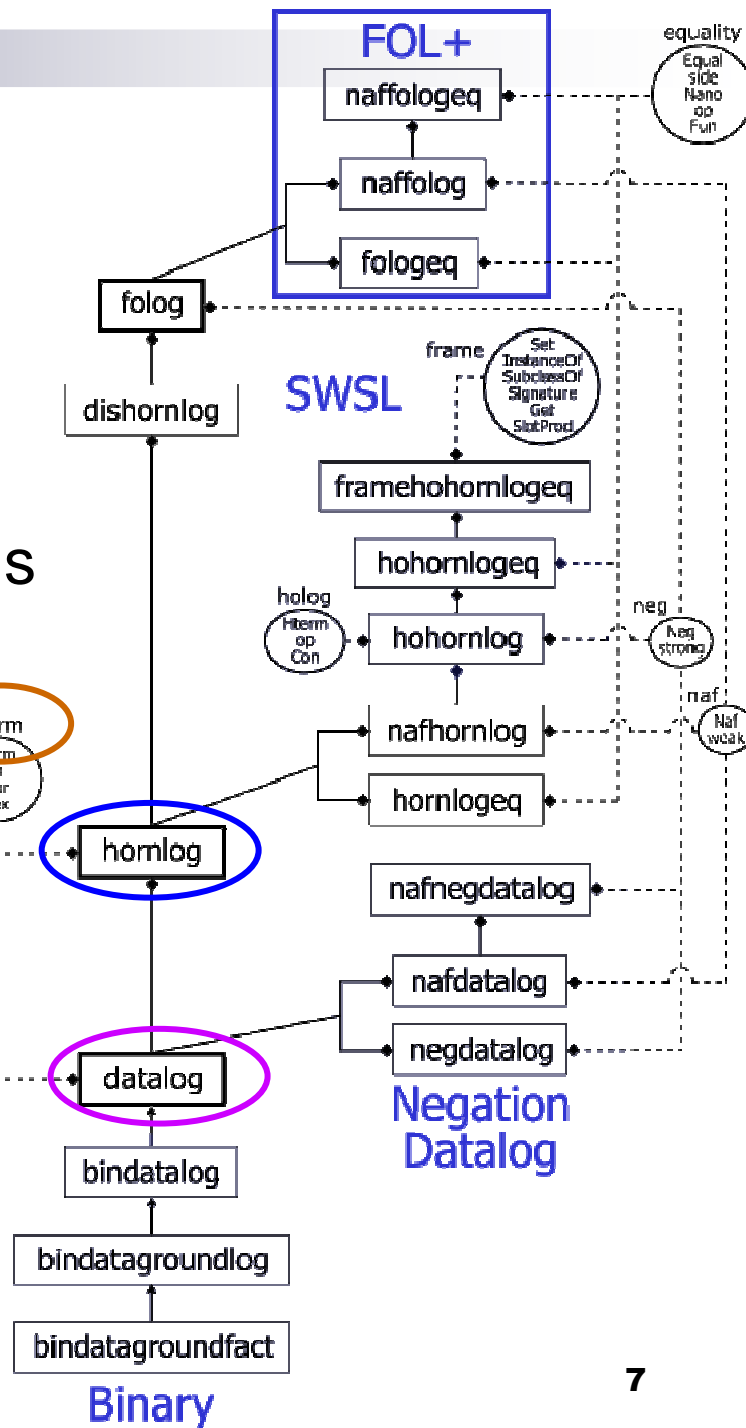
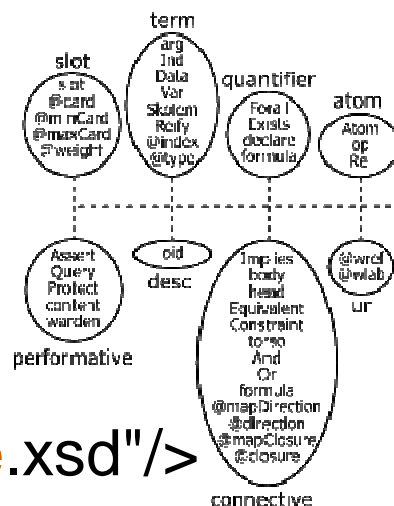
- Inheritance between schemas

- **Ovals** are auxiliary modules

- Elementary, including only element and/or attribute definitions
- Become *part of* sublanguages

E.g., in <http://www.ruleml.org/0.91/xsd/hornlog.xsd>

```
<xs:redefine
schemaLocation=
"datalog.xsd">
<xs:include
schemaLocation=
"modules/cterm_module.xsd"/>
```



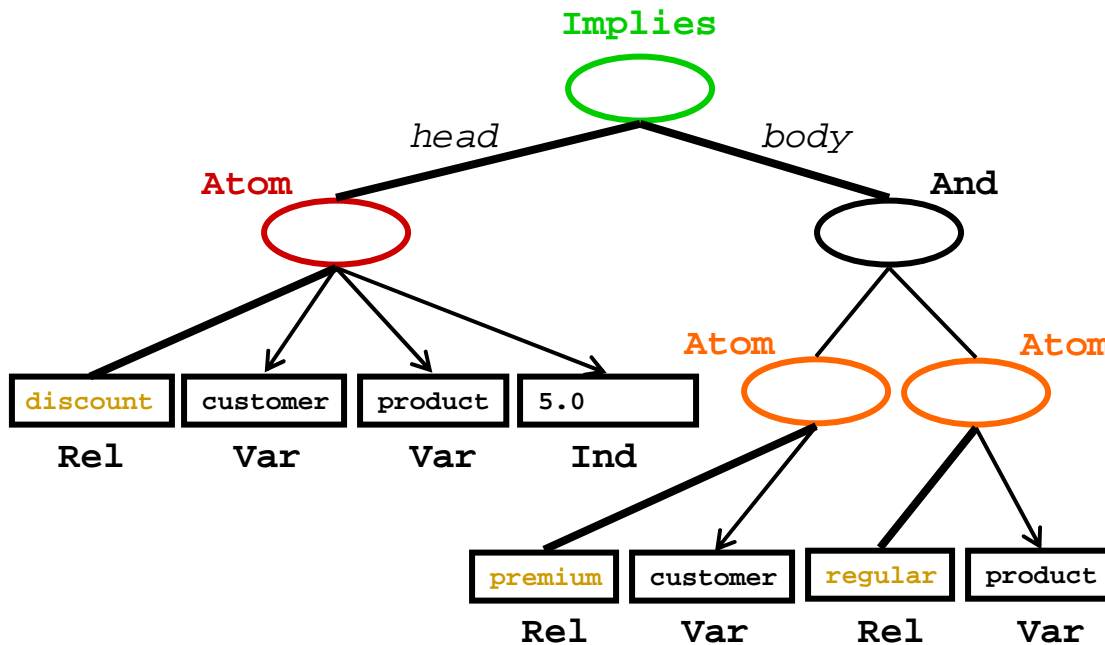


Bring Datalog to the Semantic Web

- Start with n-ary relations (not binary properties)
- Keep **V**ariable typing optional (reuse RDFS' subClassOf taxonomies as sort lattices)
- Allow signature declarations of arities and types
- Employ function-free facts as well as Horn rules (rather than 1st: RDF descriptions; 2nd: RDF rules)
- Use function-free Herbrand model semantics (querying stays decidable)
- Provide three syntactic levels:
 - User-oriented: Prolog-like, but with “?”-variables
 - Abstract: MOF/UML diagrams
 - XML serialization: Datalog RuleML

Business Rule: Positional

"The **discount** for a *customer* buying a *product* is 5 percent if the *customer* is **premium** and the *product* is **regular**."

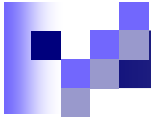


```
<Implies>
<head>
  <Atom>
    <Rel>discount</Rel>
    <Var>customer</Var>
    <Var>product</Var>
    <Ind>5.0</Ind>
  </Atom>
</head>
<body>
  <And>
    <Atom>
      <Rel>premium</Rel>
      <Var>customer</Var>
    </Atom>
    <Atom>
      <Rel>regular</Rel>
      <Var>product</Var>
    </Atom>
  </And>
</body>
</Implies>
```



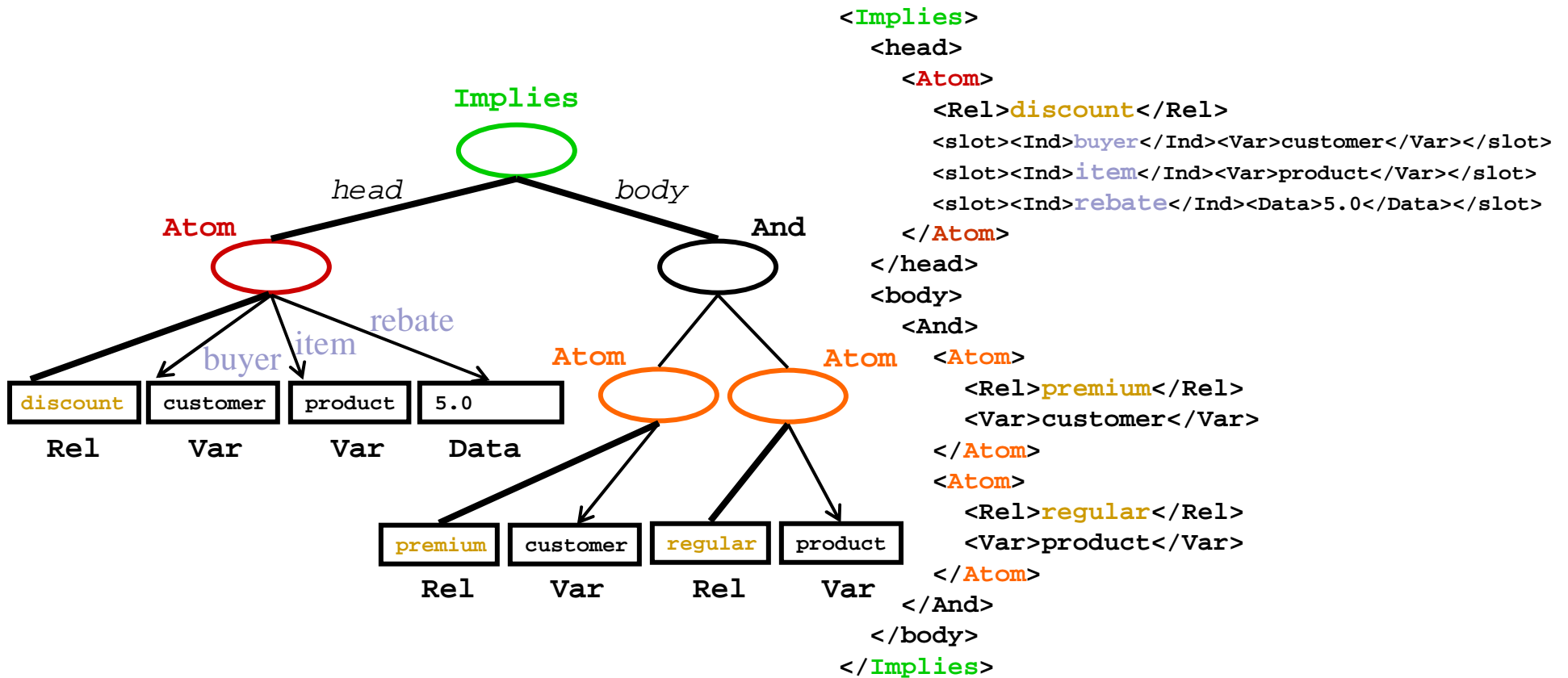
Extend Datalog for the Semantic Web (I)

- Allow slots as name- \rightarrow filler pairs in **Atoms**
(cf. F-logic's methods and RDF's properties)
- Extend optional types and signatures for slots
- Add optional object identifiers (**oids**) to atoms
- Separate **Data** literals from **Individual** constants



Business Rule: Slotted (for OO)

"The **discount** for a *customer* buying a *product* is 5 percent if the *customer* is **premium** and the *product* is **regular**."





Extend Datalog for the Semantic Web (II)

- Permit IRI webizing for **Data** (XML Schema Part 2), **Individuals** (RDF's **resources**), **Relations**, **slot** names, types (RDFS' classes), and **oids** (RDF's **about**)
- Introduce **Module** (scope) construct for clauses (cf. RDF's named graphs)
- Add scoped-default (**Naf**), strong (**Neg**), scoped-default-of-strong negation (unscoped: cf. [ERDF](#))
- Integrate with Description Logics
 - Homogeneous (SWRL, Datalog RuleML + OWL-DL)
 - Hybrid (AL-log, Datalog^{DL}, DL+log, ...)



Bring Horn Logic to the Semantic Web

- Augment Datalog with uninterpreted **F**unctions and their **E**xpressions; also for extended Datalog
- Augment Datalog's Herbrand model semantics with such **F**unctions (querying becomes undecidable)
- Extend Datalog syntaxes
 - XML Schema of Hornlog RuleML inherits and augments XML Schema of Datalog RuleML
- Add **E**quality and **i**nterpreted **F**unctions (XML serialization: attribute **in="yes"**)
- Reuse XQuery/XPath functions and operators as built-ins



Specify a First-Order Logic Web Language

- Layer on top of either
 - Disjunctive Datalog: **Or** in the head generalizing Datalog
 - Disjunctive Horn Logic: **Or** in head of near-Horn clauses
- Alternatively, layer on top of either
 - Disjunctive Datalog with restricted strong **Negation**
 - Disjunctive Horn Logic with restricted strong **Neg**
- Permit unrestricted **Or**, **And**, strong **Neg**, and quantifiers **Forall** and **Exists** to obtain FOL
- Use semantics of classical FOL model theory
- Extend Hornlog RuleML syntax to FOL RuleML



Equality for Functions

- Functional programming (FP) plays increasing Web role: MathML, XSLT, XQuery
- Functional RuleML employs orthogonal notions freely combinable with Relational RuleML
- Also solves a Relational RuleML issue, where the following ‘child-of-parent’ elements are separated:
 - Constructor (**Ctor**) of a complex term (**Cterm**)
 - User-defined function (**Fun**) of a call (**Nano**)
- Proceed to a **logic with equality**



Function Interpretedness (I)

- Different notions of ‘function’ in LP and FP:
- **LP:** *Uninterpreted functions* **denote** unspecified values when applied to arguments, not using function definitions
- **FP:** *Interpreted functions* **compute** specified returned values when applied to arguments, using function definitions
- E.g.: **first-born**: $\text{Man} \times \text{Woman} \rightarrow \text{Human}$
 - Uninterpreted: **first-born(John, Mary)** denotes first-born
 - Interpreted: using **first-born(John, Mary) = Jory**, so the application returns **Jory**



Function Interpretedness (II)

- Uninterpreted **<Ctor>** vs. interpreted **<Fun>** functions now distinguished with attribute values: **<Fun in="no">** vs. **<Fun in="yes">**
- Function applications with **Cterm** vs. **Nano** then uniformly become **Expressions**
- Two versions of example marked up as follows (where "u" stands for "no" or "yes"):

<Expr>

<Fun in="u">first-born</Fun>

<Ind>John</Ind>

<Ind>Mary</Ind>

</Expr>



Unconditional Equations

- Modified `<Equal>` element permits both symmetric and oriented equations
- E.g.: `first-born(John, Mary) = Jory` can now be marked up thus:

```
<Equal oriented="yes">  
  <lhs>  
    <Expr>  
      <Fun in="yes">first-born</Fun>  
      <Ind>John</Ind>  
      <Ind>Mary</Ind>  
    </Expr>  
  </lhs>  
  <rhs>  
    <Ind>Jory</Ind>  
  </rhs>  
</Equal>
```

Conditional Equations

- Use `<Equal>` as the conclusion of an `<Implies>`, whose condition may employ other equations
- E.g.: `?B = birth-year(?P) ⇒ age(?P) = subtract(this-year(),?B)`

`<Implies>`

`<Equal oriented="no">`

`<Var>B</Var>`

`<Expr>`

`<Fun in="yes">birth-year</Fun>`

`<Var>P</Var>`

`</Expr>`

`</Equal>`

`<Equal oriented="yes">`

`<Expr>`

`<Fun in="yes">age</Fun>`

`<Var>P</Var>`

`</Expr>`

`<Expr>`

`<Fun in="yes">subtract</Fun>`

`<Expr>`

`<Fun in="yes">this-year</Fun>`

`</Expr>`

`<Var>B</Var>`

`</Expr>`

`</Equal>`

`</Implies>`



Accommodate SWSL-Rules

- HiLog: Higher-order **V**ariables, **C**onstants, and **H**terms (complex terms and atomic formulas at the same time)
- **Equal**: As in Horn Logic with (**unoriented**) **E**quality
- Frames:
 - Value molecules: **A**toms with an **oid**, an optional **Rel** class, and zero or more name->filler instance **s**lots
 - Signature molecules: name=>filler class **s**lots, which can have {min : max} cardinality constraints
- Reification: A formula (e.g., a rule) embedded in a **Reify** element is treated (e.g., unified) as a term
- **Skolems**: Unnamed, represent new individual constants (like RDF's blank nodes); otherwise, uniquely named ones



HiLog Examples: Hterms (I)

- First-order terms: $f(a, ?X)$

<Hterm>

<op><Con>f</Con></op>

<Con>a</Con>

<Var>X</Var>

</Hterm>

- Variables over function symbols: $?X(a, ?Y)$

<Hterm>

<op><Var>X</Var></op>

<Con>a</Con>

<Var>Y</Var>

</Hterm>



HiLog Examples: Hterms (II)

- Parameterized function symbols: $f(?X, a)(b, ?X(c))$

<Hterm>

<op>

<Hterm>

<op><Con>f</Con></op>

<Var>X</Var>

<Con>a</Con>

</Hterm>

</op>

<Con>b</Con>

<Hterm>

<op><Var>X</Var></op>

<Con>c</Con>

</Hterm>

</Hterm>

Equality Example

- Equality ::= in rule head: $f(a, ?X) ::= g(?Y, b) \text{ :- } p(?X, ?Y).$

```
<Implies>
  <head>
    <Equal>
      <Hterm>
        <op><Con>f</Con></op>
        <Con>a</Con>
        <Var>X</Var>
      </Hterm>
      <Hterm>
        <op><Con>g</Con></op>
        <Var>Y</Var>
        <Con>b</Con>
      </Hterm>
    </Equal>
  </head>
  <body>
    <Hterm>
      <op><Con>p</Con></op>
      <Var>X</Var>
      <Var>Y</Var>
    </Hterm>
  </body>
</Implies>
```




Frame Example: Value Molecule

- Parameterized-name->filler slot: `o[f(a,b) -> 3]`

```
<Atom>
  <oid><Con>o</Con></oid>
  <slot>
    <Hterm>
      <op><Con>f</Con></op>
      <Con>a</Con>
      <Con>b</Con>
    </Hterm>
    <Con>3</Con>
  </slot>
</Atom>
```

Reification Example: Reified Rule

- \$Rule as slot filler: john[believes -> \${p(?X) implies q(?X)}].

```
<Hterm>
  <oid>john</oid>
  <slot>
    <Con>believes</Con>
    <Reify>
      <Implies>
        <body>
          <Hterm>
            <op><Con>p</Con></op>
            <Var>X</Var>
          </Hterm>
        </body>
      <head>
        <Hterm>
          <op><Con>q</Con></op>
          <Var>X</Var>
        </Hterm>
      </head>
    </Implies>
  </Reify>
</slot>
</Hterm>
```



Skolem Examples (I):

- Named Skolem: `holds(a, _#1)` and `between(1, _#1, 5)`.

`<And>`

`<Hterm>`

`<op><Con>holds</Con></op>`

`<Con>a</Con>`

`<Skolem>1</Skolem>`

`</Hterm>`

`<Hterm>`

`<op><Con>between</Con></op>`

`<Con>1</Con>`

`<Skolem>1</Skolem>`

`<Con>5</Con>`

`</Hterm>`

`</And>`



Skolem Examples (II):

- Unnamed Skolem: `holds(a, _#)` and `between(1, _#, 5)`.

`<And>`

`<Hterm>`

`<op><Con>holds</Con></op>`

`<Con>a</Con>`

`<Skolem/>`

`</Hterm>`

`<Hterm>`

`<op><Con>between</Con></op>`

`<Con>1</Con>`

`<Skolem/>`

`<Con>5</Con>`

`</Hterm>`

`</And>`



Proceed towards Modal Logics

- Modal operators **generically** viewed as special **Re**lations at least one of whose arguments is a proposition represented as an **Atom** with an uninterpreted **Re**lation (including another modal operator, but not an arbitrary formula)
 - *Alethic* **necessary** (\square) and **possible** (\diamond)
 - *Deontic* **must** and **may** (e.g., in business rules)
 - Open for *temporal* (e.g., when planning/diagnosing reactive rules), *epistemic* (e.g., in authentication rules), and further modal operators
- Towards a [unified framework](#) for multi-modal logic based on Kripke-style possible worlds semantics



Modal Examples: Alethic Operator

- Necessity: $\Box \text{prime}(1)$

```
<Atom>
```

```
  <Rel modal="yes">necessary</Rel>
```

```
  <Atom>
```

```
    <Rel in="no">prime</Rel>
```

```
    <Data>1</Data>
```

```
  </Atom>
```

```
</Atom>
```



Modal Examples: Epistemic Operator

- Knowledge: `knows(Mary,material(moon,rock))`

`<Atom>`

`<Rel modal="yes">knows</Rel>`

`<Ind>Mary</Ind>`

`<Atom>`

`<Rel in="no">material</Rel>`

`<Ind>moon</Ind>`

`<Ind>rock</Ind>`

`</Atom>`

`</Atom>`



Modal Examples: Epistemic Reasoning

- Veridicality axiom: $Knows_{Agent} proposition \rightarrow proposition$

$Knows_{Mary} material(moon, rock) \rightarrow material(moon, rock)$

Serialization in previous slide

→

<Atom>

<Rel in="yes">material</Rel> <!-- "yes" is default -->

<Ind>moon</Ind>

<Ind>rock</Ind>

</Atom>



Modal Examples: Nested Operators

- Knowledge of Necessity: `knows (Mary, □ prime (1))`

```
<Atom>
```

```
  <Rel modal="yes">knows</Rel>
```

```
  <Ind>Mary</Ind>
```

```
  <Atom>
```

```
    <Rel modal="yes" in="no">necessary</Rel>
```

```
    <Atom>
```

```
      <Rel in="no">prime</Rel>
```

```
      <Data>1</Data>
```

```
    </Atom>
```

```
  </Atom>
```

```
</Atom>
```



Protect Knowledge Bases by Integrity Constraints

- A knowledge base KB is a formula in any of our logic languages
- An integrity constraint IC is also a formula in any of our logic languages, which may be chosen independently from KB
- KB obeys IC
iff
KB entails IC
(Reiter [1984](#), [1987](#))
 - Entailment notion of [1987](#) uses epistemic modal operator
- Serialization: **<Entails> KB IC </ Entails>**

Integrity Constraint Example: Rule with \exists -Head

- Adapted from (Reiter [1987](#)):

$$IC = \{ (\forall \mathbf{x}) \text{emp}(\mathbf{x}) \Rightarrow (\exists \mathbf{y}) \text{ssn}(\mathbf{x}, \mathbf{y}) \}$$

$$KB_1 = \{ \text{emp}(\text{Mary}) \}$$

KB₁ violates IC

$$KB_2 = \{ \text{emp}(\text{Mary}), \text{ssn}(\text{Mary}, 1223) \}$$

KB₂ obeys IC

<Entails> KB_i IC </Entails>

KB₁:

```
<Atom>
  <Rel>emp</Rel>
  <Ind>Mary</Ind>
</Atom>
```

KB₂:

```
<Rulebase>
  <Atom>
    <Rel>emp</Rel>
    <Ind>Mary</Ind>
  </Atom>
  <Atom>
    <Rel>ssn</Rel>
    <Ind>Mary</Ind>
    <Data>1223</Data>
  </Atom>
</Rulebase>
```

IC:

```
<Forall>
  <Var>x</Var>
  <Implies>
    <Atom>
      <Rel>emp</Rel>
      <Var>x</Var>
    </Atom>
    <Exists>
      <Var>y</Var>
      <Atom>
        <Rel>ssn</Rel>
        <Var>x</Var>
        <Var>y</Var>
      </Atom>
    </Exists>
  </Implies>
</Forall>
```



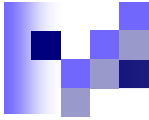
Approach Production and Reaction Rules

- Share Condition (C) part with earlier languages as proposed for the [RIF Condition Language](#)
- Develop Action (A) part of Production Rules via a taxonomy of actions on KBs (Assert, Retract, ...), on local or remote hosts, or on the surroundings
- Develop Event (E) part of Reaction Rules via a corresponding taxonomy
- Create CA and ECA families bottom-up and map to relevant languages for Semantic Web Services
- Serialized: **<Reaction> E C A </Reaction>**
- See <http://ibis.in.tum.de/research/ReactionRuleML> TG



RDF Rules

- RDF-like Rules: Important RuleML sublanguage
 - Datalog: Relational databases augmented by views
 - RDF Properties: Slots permit non-positional, keyed arguments
 - RDF URIs/IRIs: Anchors provide **object identity** via webzing through URIs/IRIs
 - **oids**: Can be **I**ndividuals, **V**ariables, etc.
 - **iris**: Now used for both RDF's **a**bout and **r**esource
 - RDF Blank Nodes: F-logic/Flora-2 Skolem-constant approach
 - E.g., Skolem generator '_' becomes `<skolem/>`



```
<Implies>
  <body>
    <And>
      <Atom>
        <oid><Var>x</Var></oid>
        <Rel>product</Rel>
        <slot><Ind iri=":price"/><Var>y</Var></slot>
        <slot><Ind iri=":weight"/><Var>z</Var></slot>
      </Atom>
      <Atom>
        <Rel iri="swrlb:greaterThan"/><Var>y</Var><Data>200</Data>
      </Atom>
      <Atom>
        <Rel iri="swrlb:lessThan"/><Var>z</Var><Data>50</Data>
      </Atom>
    </And>
  </body>
  <head>
    <Atom>
      <oid><Var>x</Var></oid>
      <Rel>product</Rel>
      <slot><Ind iri=":shipping"/><Data>0</Data></slot>
    </Atom>
  </head>
</Implies>
```

“For a product whose price is greater than 200 and whose weight is less than 50, no shipping is billed.”



Bidirectional Interpreters in Java

- Two varieties of reasoning engines
 - **Top-Down**: backward chaining
 - **Bottom-Up**: forward chaining
- **jDREW**: *Java Deductive Reasoning Engine for the Web* includes both TD and BU
<http://www.jdrew.org>
- **OO jDREW**: *Object-Oriented* extension to jDREW
<http://www.jdrew.org/oojdrew>
- Java Web Start online demo available at
<http://www.jdrew.org/oojdrew/demo.html>



OO jDREW Slots

- Normalized atoms and complex terms

- **oids** (object identifier)
- **Positional** parameters (in their original order)
- **Positional** rest terms
- **Slotted** parameters (in the order encountered)
- **Slotted** rest terms

- Efficient unification algorithm

- Linear $O(m+n)$: instead of $O(m*n)$
 - No need for positional order
 - Slots internally sorted
- Steps:
 - Scan two lists of parameters
 - Matching up roles and positions for positional parameters
 - Unifying those parameters
 - Add unmatched roles to list of rest terms
 - Generate dynamically a Plex (RuleML's closest equivalent to a list) for a collection of rest terms

The screenshot shows the OO jDREW Top-Down Engine interface. The 'Query' field contains the text: `discount(?person, ?thing, ?amount).` Below the query field are buttons for 'Issue Query' and 'Next Solution'. The 'Solution' field displays a tree view of the solution: `$top() :- discount(PeterMiller, Honda, percent5).` with sub-entries for `discount(PeterMiller, Honda, percent5) :- premium(PeterMiller), regular(Honda).` and `premium(PeterMiller).` To the right, the 'Variable Bindings' table shows the following data:

Variable	Binding
?person	PeterMiller
?thing	Honda
?amount	percent5

At the bottom of the window is a 'Show Debug Console' button. The status bar at the very bottom reads 'Java Application Window'.

positional

POSL
syntax

**discount(?customer,?product,percent5)
:- premium(?customer), regular(?product).

premium(PeterMiller).
regular(Honda).**

The screenshot shows the OO jDREW Top-Down Engine interface. The 'Query' tab is active, displaying the query: `discount(rebate->?amount;prod->?thing;cust->?person).` Below the query, there are buttons for 'Issue Query' and 'Next Solution'. The 'Solution' section shows a tree view of the solution: `$top() :- discount(cust->PeterMiller; prod->Honda; rebate->percent5).` with sub-nodes for `premium(cust->PeterMiller).` and `regular(prod->Honda).` To the right, the 'Variable Bindings' table is shown:

Variable	Binding
?person	PeterMiller
?thing	Honda
?amount	percent5

At the bottom right, there is a 'Show Debug Console' button.



slotted

POSL
syntax

**discount(cust->?customer;prod->?product;rebate->percent5)
:- premium(cust->?customer), regular(prod->?product).**

**premium(cust->PeterMiller).
regular(prod->Honda).**



OO jDREW Types

- Order-sorted type system
 - RDF Schema: lightweight taxonomies of the Semantic Web
 - To specify a partial order for a set of classes in RDFS
- Advantages
 - Having the appropriate types specified for the parameters
 - To restrict the search space
 - Faster and more robust system than when reducing types to unary predicate calls in the body
- Limitations
 - Only modeling the taxonomic relationships between classes
 - Not modeling properties with domain and range restrictions

OO jDREW Top-Down Engine

Type Definition Knowledge Base Query

Query:
`base_price(customer->[sex->male; name->"John Doe"; age->28]; vehicle->vehicle:ToyotaCorolla; price->?money:Integer).`

Issue Query Next Solution

Solution:
`$top():-base_price(customer->[sex->male; name->"John Doe"; age->28]; vehicle->:Car; price->650:Integer).`
`base_price(customer->[sex->male; name->"John Doe"; age->28!?!]; vehicle->:Van; price->725:Integer).`

Variable Bindings:

Variable	Binding
?money : Integer	650 : Integer

```

classDiagram
    class Thing
    class Vehicle
    class PassengerVehicle
    class Van
    class Car
    class MiniVan
    class Sedan
    class ToyotaCorolla
    class Nothing

    Thing <|-- Vehicle
    Vehicle <|-- PassengerVehicle
    Vehicle <|-- Van
    PassengerVehicle <|-- Car
    PassengerVehicle <|-- MiniVan
    Car <|-- Sedan
    Car <|-- Sedan
    ToyotaCorolla <|-- Sedan
    Nothing ..> ToyotaCorolla
    Nothing ..> Sedan
    Nothing ..> MiniVan
  
```

`base_price(customer->[sex->male;!?!]; vehicle->:Car; price->650:Integer).`

`base_price(customer->[sex->male;!?!]; vehicle->:Van; price->725:Integer).`

Java Application Window

OO jDREW OIDs

- `oid`: Object Identifier
- Currently: symbolic names
 - In `<Atom>` & `<Implies>`
- Planned: `iri` attribute
- E.g., give name to fact `keep(Mary, ?object)`.

```
<Atom>
  <oid><Ind>mary-12</Ind></oid>
  <Rel>keep</Rel>
  <Ind>Mary</Ind>
  <Var>object</Var>
</Atom>

<Atom>
  <oid><Ind iri="http://mkb.ca"/></oid>
  <Rel>keep</Rel>
  <Ind>Mary</Ind>
  <Var>object</Var>
</Atom>

<Atom>
  <oid><Var>object</Var></oid>
  <Rel>keep</Rel>
  <Ind>Mary</Ind>
  <Var>object</Var>
</Atom>
```



Conclusions

- [RuleML](#) is modular family, whose root allows to access the language as a whole and whose members allow customized subsets
- New members joining, e.g. [Fuzzy RuleML](#)
- Concrete & abstract syntax of RuleML
 - Specified by modular XSD (shown here) & MOF
- Formal semantics of OO Hornlog RuleML
 - Implemented by OO jDREW BU & TD
- Interoperability/Interchange of/with RuleML
 - Realized by translators, primarily via XSLT